# Discrete Mathematics 

## Lecture 03

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## Negating Quantified Expressions (1/4)

## كلية الحاسبات والذكاء الإصطناعي

## Negating Quantified Expressions:

$P(x)$ is the statement " $x$ has taken a course in calculus" and the domain consists of the students in your class.
$\forall \boldsymbol{x P}(x):$

## Negating Quantified Expressions (1/4)

## كلية الحاسبات والذكاء الإصطناعي

## Negating Quantified Expressions:

$P(x)$ is the statement " $x$ has taken a course in calculus" and the domain consists of the students in your class.
$\forall \boldsymbol{x P}(x)$ :
"Every student in your class has taken a course in calculus"

## Negating Quantified Expressions (2/4)

## كلية الحاسبات والذكاء الإصطناعي

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"Every student in your class has taken a course in calculus"
The negation of this statement is
"There is at least one student in your class who has not taken a course in calculus"

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$$
\neg \forall \boldsymbol{x P}(\boldsymbol{x})
$$

## Negating Quantified Expressions (2/4)

## كلية الحاسبات والذكاء الإصطناعي

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"Every student in your class has taken a course in calculus"
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$$
\neg \forall x P(x) \equiv \exists x \neg P(x)
$$

## Negating Quantified Expressions (3/4)

## كلية الحاسبات والذكاء الإصطناعي

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## Negating Quantified Expressions (4/4)

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$P(x)$ is the statement " $x$ has taken a course in calculus" and the domain consists of the students in your class.
$\exists x P(x):$
"At least one student in your class has taken a course in calculus"
The negation of this statement is
"Every student in this class has not taken calculus"

## Negating Quantified Expressions (4/4)

## كلية الحاسبات والذكاء الإصطناعي

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$\exists x P(x):$
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$$
\neg \exists \boldsymbol{x P}(x)
$$

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## كلية الحاسبات والذكاء الإصطناعي

## Negating Quantified Expressions:

$P(x)$ is the statement " $x$ has taken a course in calculus" and the domain consists of the students in your class.
$\exists x P(x):$
"At least one student in your class has taken a course in calculus"
The negation of this statement is
"Every student in this class has not taken calculus"

$$
\neg \exists x P(x) \equiv \forall x \neg P(x)
$$

## Rules of Inference (1/9)

## Valid Arguments in Propositional Logic

Consider the following argument involving propositions (which, by definition, is a sequence of propositions):
"If you have a current password, then you can log onto the network."
"You have a current password."
Therefore,
"You can $\log$ onto the network."

## Rules of Inference (1/9)

## Valid Arguments in Propositional Logic

Consider the following argument involving propositions (which, by definition, is a sequence of propositions):
"If you have a current password, then you can log onto the network."
"You have a current password."
Premises
Therefore,
"You can $\log$ onto the network."
Conclusion

## Rules of Inference (1/9)

## Valid Arguments in Propositional Logic

Consider the following argument involving propositions (which, by definition, is a sequence of propositions):

$$
\begin{aligned}
& p \rightarrow q \\
& p
\end{aligned}
$$

## Premises

$\therefore q$
Conclusion

## Rules of Inference (1/9)

## Valid Arguments in Propositional Logic

Consider the following argument involving propositions (which, by definition, is a sequence of propositions):

$$
\begin{aligned}
& p \rightarrow q \\
& p
\end{aligned}
$$

## Premises

$\therefore q$
Conclusion
This argument is valid if $((p \rightarrow q) \wedge p) \rightarrow q$ is a tautology.

## Rules of Inference (1/9)

## Valid Arguments in Propositional Logic

An argument in propositional logic is a sequence of propositions. All the proposition in the argument are called premises and the final proposition is called the conclusion.

$$
\begin{aligned}
& p \rightarrow q \\
& p
\end{aligned}
$$

## Premises

$\therefore q$
Conclusion
This argument is valid if $((p \rightarrow q) \wedge p) \rightarrow q$ is a tautology.

## Rules of Inference (2/9)

## Valid Arguments in Propositional Logic

We can always use a truth table to show that an argument form is valid.

$$
p \rightarrow q
$$

$p$
$\therefore q$

## Rules of Inference (2/9)

## Valid Arguments in Propositional Logic

We can always use a truth table to show that an argument form is valid.

Premise 1

$p$
$\therefore q$

## Rules of Inference (2/9)

## Valid Arguments in Propositional Logic

We can always use a truth table to show that an argument form is valid.

Premise 1 Premise 2


$p$
$\therefore q$

## Rules of Inference (2/9)

## Valid Arguments in Propositional Logic

We can always use a truth table to show that an argument form is valid.

Premise 1 Premise 2

|  | $p$ | $q$ | $p \rightarrow q$ | $p$ | $(p \rightarrow q) \wedge p$ |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  |  |  |  |  |  |  |
|  | $\mathbf{T}$ | $\mathbf{T}$ | $\mathbf{T}$ | $\mathbf{T}$ | $\mathbf{T}$ |  |
|  | $\mathbf{T}$ | $\mathbf{F}$ | $\mathbf{F}$ | $\mathbf{T}$ | $\mathbf{F}$ |  |
|  | $\mathbf{F}$ | $\mathbf{T}$ | $\mathbf{T}$ | $\mathbf{F}$ | $\mathbf{F}$ |  |
| $\boldsymbol{p} \rightarrow \boldsymbol{q}$ | $\mathbf{F}$ | $\mathbf{F}$ | $\mathbf{T}$ | $\mathbf{F}$ | $\mathbf{F}$ |  |

$p$
$\therefore q$

## Rules of Inference (2/9)

## Valid Arguments in Propositional Logic

We can always use a truth table to show that an argument form is valid.

$$
\text { Premise } 1 \text { Premise } 2 \quad \text { Conclusion }
$$

|  | $p$ | $q$ | $p \rightarrow q$ | $p$ | $(p \rightarrow q) \wedge p$ | $q$ |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | T | T | T | T | T | T |  |
|  | T | F | F | T | F | F |  |
|  | F | T | T | F | F | T |  |
| $p \rightarrow q$ | F | F | T | F | F | F |  |

$p$
$\therefore q$

## Rules of Inference (2/9)

## Valid Arguments in Propositional Logic

We can always use a truth table to show that an argument form is valid.

$$
\text { Premise } 1 \text { Premise } 2 \quad \text { Conclusion }
$$

|  | $p$ | $q$ | $p \rightarrow q$ | $p$ | $(p \rightarrow q) \wedge p$ | $q$ | $((p \rightarrow q) \wedge p) \rightarrow q$ |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | T | T | T | T | T | T | T |
|  | T | F | F | T | F | F | T |
|  | F | T | T | F | F | T | T |
| $p \rightarrow q$ | F | F | T | F | F | F | T |

$p$

$$
((p \rightarrow q) \wedge p) \rightarrow q \text { is a tautology }
$$

$\therefore q$

## Rules of Inference (3/9)

## TABLE 1 Rules of Inference.

Part 1

| Rule of Inference | Tautology | Name |
| :---: | :--- | :--- |
| $p$ | $(p \wedge(p \rightarrow q)) \rightarrow q$ | Modus ponens |
| $\therefore \frac{p \rightarrow q}{q}$ |  |  |
| $\neg q$ | $(\neg q \wedge(p \rightarrow q)) \rightarrow \neg p$ | Modus tollens |
| $\therefore \frac{p \rightarrow q}{\neg p}$ |  |  |
| $p \rightarrow q$ <br> $q \rightarrow r$ | $((p \rightarrow q) \wedge(q \rightarrow r)) \rightarrow(p \rightarrow r)$ | Hypothetical syllogism |
| $\therefore \frac{p \rightarrow r}{}$ |  |  |
| $\quad p \vee q$ | $((p \vee q) \wedge \neg p) \rightarrow q$ | Disjunctive syllogism |
| $\therefore \frac{\neg p}{q}$ |  |  |

## Rules of Inference (3/9)

TABLE 1 Rules of Inference.
Part 2

| Rule of Inference | Tautology | Name |
| :---: | :--- | :--- |
| $\therefore \overline{p \vee q}$ | $p \rightarrow(p \vee q)$ | Addition |
| $\quad$$p \wedge q$ <br> $p$ | $(p \wedge q) \rightarrow p$ | Simplification |
| $p$ | $((p) \wedge(q)) \rightarrow(p \wedge q)$ | Conjunction |
| $\therefore \frac{q}{p \wedge q}$ |  |  |
| $p \vee q$ | $((p \vee q) \wedge(\neg p \vee r)) \rightarrow(q \vee r)$ | Resolution |
| $\therefore \frac{\neg p \vee r}{q \vee r}$ |  |  |

## Rules of Inference (4/9)

## كلية الحاسبات والذكاء الإصطناعي

## Example1

Using the truth table to show that the hypotheses
$p \vee q$
$\neg p \vee r$
lead to the conclusion
$q \vee r$

|  | $p \vee q$ | $((p \vee q) \wedge(\neg p \vee r)) \rightarrow(q \vee r)$ |
| :--- | :--- | :--- |
|  | $\neg p \vee r$ | Resolution |
| $\therefore$ | $q \vee r$ |  |

## Rules of Inference (4/9)

## كلية الحاسبات والذكاء الإصطناعي

## Example1

Using the truth table to show that the hypotheses

| $p \vee q$ | Premise 1 |  |  |  |  | Premise 2 C |  | onclusion |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | $p$ | $q$ | $r$ | $p \vee q$ | $\neg p$ | $\neg p \vee r$ | $(p \vee q) \wedge(\neg p \vee r)$ | $q \vee r$ |
| $\neg p \vee r$ | T | T | T | T | F | T | T | T |
|  | T | T | F | T | F | F | F | T |
| $q \vee r$ | T | F | T | T | F | T | T |  |
|  | T | F | F | T | F | F | F | F |
|  | F | T | T | T | T | T | T | T |
|  | F | T | F | T | T | T | T |  |
|  | F | F | T | F | T | T | F | T |
|  | F | F | F | F | T | T | F | F |

## Rules of Inference (5/9)

## كلية الحاسبات والذكاء الإصطناعي

## Example2

Using the rules of inference to show that the hypotheses
$\neg p \wedge q$
$r \rightarrow p$
$\neg r \rightarrow S$
$s \rightarrow t$
lead to the conclusion
$t$

## Rules of Inference (5/9)

## Example2

$\neg p \wedge q$
$\therefore \neg p$

| $\therefore p \wedge q$ | $(p \wedge q) \rightarrow p$ | Simplification |
| :--- | :--- | :--- |

$$
\begin{aligned}
& \neg p \wedge q \\
& r \rightarrow p \\
& \neg r \rightarrow s \\
& s \rightarrow t
\end{aligned}
$$

## Rules of Inference (5/9)

## كلية الحاسبات والذكاء الإصطناعي

## Example2

$\neg p \wedge q$
$\therefore \neg p$

| $\quad p \wedge q$ | $(p \wedge q) \rightarrow p$ | Simplification |
| :--- | :--- | :--- |
| $\therefore \bar{p}$ |  |  |


| $\begin{aligned} & \neg p \\ & r \rightarrow p \end{aligned}$ | $\begin{aligned} & \neg q \\ & \therefore p \rightarrow q \\ & \neg p \end{aligned}$ | $(\neg q \wedge(p \rightarrow q)) \rightarrow \neg p$ | Modus tollens |
| :---: | :---: | :---: | :---: |
| $\therefore \neg r$ |  |  |  |

## Example2

$\neg r$
$\neg r \rightarrow s$

| $p$ | $(p \wedge(p \rightarrow q)) \rightarrow q$ | Modus ponens |
| :---: | :--- | :--- |
| $\therefore \frac{p \rightarrow q}{q}$ |  |  |

$\therefore S$


## Rules of Inference (5/9)

## كلية الحاسبات والذكاء الإصطناعي

## Example2

$\neg r$
$\neg r \rightarrow S$

| $p$ | $(p \wedge(p \rightarrow q)) \rightarrow q$ | Modus ponens |
| :---: | :--- | :--- |
| $\therefore \frac{p \rightarrow q}{q}$ |  |  |

$\therefore S$
$S$
$s \rightarrow t$
$\therefore t$
conclusion


## Chapter 2: Basic Structures

## كلية الحاسبات والذكاء الإصطناعي

- Sets.
- Functions.
- Sequences, and Summations.
- Matrices.


## Sets (1/24)

A set is an unordered collection of objects.

The objects in a set are called the elements, or members, of the set. A set is said to contain its elements.

## Sets (2/24)

## كلية الحاسبات والذكاء الإصطناعي

$S=\{a, b, c, d\}$
We write $a \in S$ to denote that $a$ is an element of the set $S$. The notation $e \notin S$ denotes that $e$ is not an element of the set $S$.

## Sets (3/24)

## كلية الحاسبات والذكاء الإصطناعي

## The set $O$ of odd positive integers less than 10 can be expressed by $O=\{1,3,5,7,9\}$.

The set of positive integers less than 100 can be denoted by $\{1,2,3, \ldots, 99\}$.
ellipses (...)

## Sets (4/24)

## Another way to describe a set is to use set builder notation.

The set $O$ of odd positive integers less than 10 can be expressed by $O=\{1,3,5,7,9\}$.
$O=\{x \mid x$ is an odd positive integer less than 10$\}$,
$O=\left\{x \in \mathbf{Z}^{+} \mid x\right.$ is odd and $\left.x<10\right\}$.
$\mathbf{N}=\{0,1,2,3, \ldots\}$, the set of all natural numbers
$\mathbf{Z}=\{\ldots,-2,-1,0,1,2, \ldots\}$, the set of all integers
$\mathbf{Z}^{+}=\{1,2,3, \ldots\}$, the set of all positive integers
$\mathbf{Q}=\{p / q \mid p \in \mathbf{Z}, q \in \mathbf{Z}$, and $q \neq 0\}$,
the set of all rational numbers
$\mathbf{R}$, the set of all real numbers
$\mathbf{R}^{+}$, the set of all positive real numbers
$\mathbf{C}$, the set of all complex numbers.

## Interval Notation

Closed interval $[a, b]$
Open interval $(a, b)$
$[a, b]=\{x \mid a \leq x \leq b\}$
$[a, b)=\{x \mid a \leq x<b\}$
$(a, b]=\{x \mid a<x \leq b\}$
$(a, b)=\{x \mid a<x<b\}$

## Sets (7/24)

If $A$ and $B$ are sets, then $A$ and $B$ are equal if and only if $\forall x(x \in A \leftrightarrow x \in B)$. We write $A=B$, if $A$ and $B$ are equal sets.

- The sets $\{1,3,5\}$ and $\{3,5,1\}$ are equal, because they have the same elements.
- $\{1,3,3,5,5,5\}$ is the same as the set $\{1,3,5\}$ because they have the same elements.


## Sets (8/24)

## Empty Set

There is a special set that has no elements. This set is called the empty set, or null set, and is denoted by $\emptyset$. The empty set can also be denoted by $\}$

## Sets (9/24)

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كلية الحاسبات والذكاء الإصطناعي
```


## Cardinality

The cardinality is the number of distinct elements in $S$. The cardinality of $S$ is denoted by $|S|$.

## Sets (10/24)

## Example1

$$
\begin{aligned}
& S=\{a, b, c, d\} \\
& |S|=4 \\
& A=\{1,2,3,7,9\} \\
& \emptyset=\{ \}
\end{aligned}
$$

## Sets (10/24)

## Example1

$$
\begin{aligned}
& S=\{a, b, c, d\} \\
& |S|=4 \\
& A=\{1,2,3,7,9\} \\
& |A|=5 \\
& \emptyset=\{ \} \\
& |\varnothing|=0
\end{aligned}
$$

## Example2

$S=\{a, b, c, d,\{2\}\}$
$|S|=$
$A=\{1,2,3,\{2,3\}, 9\}$
$|A|=$
$\{\varnothing\}=\{\{ \}\}$
$|\{\varnothing\}|=$

## Example2

$S=\{a, b, c, d,\{2\}\}$
$|S|=5$
$A=\{1,2,3,\{2,3\}, 9\}$
$|A|=5$
$\{\varnothing\}=\{\{ \}\}$
$|\{\varnothing\}|=1$

## Sets (12/24)

## Infinite

A set is said to be infinite if it is not finite.
The set of positive integers is infinite.

$$
Z^{+}=\{1,2,3, \ldots\}
$$

## Sets (13/24)

## Subset

The set $A$ is said to be a subset of $B$ if and only if every element of $A$ is also an element of $B$.

We use the notation $A \subseteq B$ to indicate that $A$ is a subset of the set $B$.

$$
A \subseteq B \leftrightarrow \forall x(x \in A \rightarrow x \in B)
$$

## Sets (13/24)

## Subset

The set $A$ is said to be a subset of $B$ if and only if every element of $A$ is also an element of $B$.

We use the notation $A \subseteq B$ to indicate that $A$ is a subset of the set $B$.

$$
(A \subseteq B) \equiv(B \supseteq A)
$$

$$
A \subseteq B \leftrightarrow \forall x(x \in A \rightarrow x \in B)
$$

## Sets (13/24)

## كلية الحاسبات والذكاء الإصطناعي

## Subset

For every set $S$,
(i) $\emptyset \subseteq S$ and (ii) $S \subseteq S$.

To show that two sets $A$ and $B$ are equal, show that $A \subseteq B$ and $B \subseteq A$.

## Sets (14/24)

## كلية الحاسبات والذكاء الإصطناعي

## Proper Subset

The set $A$ is a subset of the set $B$ but that $A \neq B$, we write $A \subset B$ and say that $A$ is a proper subset of $B$.

$$
A \subset B \leftrightarrow(\forall x(x \in A \rightarrow x \in B) \wedge \exists x(x \in B \wedge x \notin A))
$$

## Sets (15/24)

## Example

For each of the following sets, determine whether 3 is an element of that set.
$\{1,2,3,4\}$
$\{\{1\},\{2\},\{3\},\{4\}\}$
$\{1,2,\{1,3\}\}$

## Sets (15/24)

## Example

For each of the following sets, determine whether 3 is an element of that set.
$3 \in\{1,2,3,4\}$
$3 \notin\{\{1\},\{2\},\{3\},\{4\}\}$
$3 \notin\{1,2,\{1,3\}\}$

## Sets (16/24)

## Venn Diagram

$$
\begin{aligned}
& A=\{1,2,3,4,7\} \\
& B=\{0,3,5,7,9\} \\
& C=\{1,2\}
\end{aligned}
$$

## Sets (17/24)

## Venn Diagram

$$
\begin{aligned}
& A=\{1,2,3,4,7\} \\
& B=\{0,3,5,7,9\} \\
& C=\{1,2\}
\end{aligned}
$$

## Sets (18/24)

## Power Set

## The set of all subsets.

If the set is $S$. The power set of $S$ is denoted by $P(S)$.
The number of elements in the power set is $2^{|S|}$

## Sets (18/24)

## Power Set

## The set of all subsets.

If the set is $S$. The power set of $S$ is denoted by $P(S)$.
The number of elements in the power set is $2^{|S|}$
$S=\{1,2,3\}$

$$
|P(S)|=2^{3}=8 \text { elements }
$$

$P(S)=2^{S}$
$=\{\emptyset,\{1\},\{2\},\{3\},\{1,2\},\{1,3\},\{2,3\},\{1,2,3\}\}$

## Sets (19/24)

## كلية الحاسبات والذكاء الإصطناعي

## Example1

## What is the power set of the empty set?

## Sets (19/24)

## كلية الحاسبات والذكاء الإصطناعي

## Example1

## What is the power set of the empty set?

$\mathcal{P}(\emptyset)=\{\emptyset\}$.

## Example2

## What is the power set of the set $\{\emptyset\}$ ?

## Example2

## What is the power set of the set $\{\emptyset\}$ ?

$$
\mathcal{P}(\{\emptyset\})=\{\emptyset,\{\emptyset\}\} .
$$

## Sets (21/24)

## The ordered $n$-tuple

The ordered $n$-tuple $\left(a_{1}, a_{2}, \ldots, a_{n}\right)$ is the ordered collection that has $a_{1}$ as its first element, $a_{2}$ as its second element, $\ldots$, and $a_{n}$ as its $n$th element.

In particular, ordered 2 -tuples are called ordered pairs (e.g., the ordered pairs $(a, b))$

## Sets (22/24)

## كلية الحاسبات والذكاء الإصطناعي

## Cartesian Products

Let $A$ and $B$ be sets.
The Cartesian product of $A$ and $B$, denoted by $A \times B$, is the set of all ordered pairs $(a, b)$, where $a \in A$ and
$b \in B$. Hence, $A \times B=\{(a, b) \mid a \in A \wedge b \in B\}$.

## Sets (22/24)

## Cartesian Products - Example

Let $A=\{1,2\}$, and $B=\{a, b, c\}$
$A \times B=\{(1, a),(1, b),(1, c),(2, a),(2, b),(2, c)\}$.

$$
|A \times B|=|A| *|B|=2 * 3=6
$$

## Sets (22/24)

## Cartesian Products - Example

Let $A=\{1,2\}$, and $B=\{a, b, c\}$
$A \times B=\{(1, a),(1, b),(1, c),(2, a),(2, b),(2, c)\}$.
$|A \times B|=|A| *|B|=2 * 3=6$

Find $B \times A$ ?

## Sets (23/24)

## The Cartesian product of more than two sets.

The Cartesian product of the sets $A_{1}, A_{2}, \ldots, A_{n}$, denoted by $A_{1} \times A_{2} \times \cdots \times A_{n}$, is the set of ordered n-tuples $\left(a_{1}, a_{2}, \ldots, a_{n}\right)$, where $a_{i}$ belongs to $A_{i}$ for $i=1,2, \ldots, n$. In other words,

$$
\begin{gathered}
A_{1} \times A_{2} \times \cdots \times A_{n}= \\
\left\{\left(a_{1}, a_{2}, \ldots, a_{n}\right) \mid a_{i} \in A_{i} \text { for } i=1,2, \ldots, n\right\} .
\end{gathered}
$$

## Sets (24/24)

## كلية الحاسبات والذكاء الإصطناعي

## Example:

$A \times B \times C$, where $A=\{0,1\}, B=\{1,2\}$, and $C=\{0,1,2\}$

$$
\begin{aligned}
A \times B \times C= & \{(0,1,0),(0,1,1),(0,1,2),(0,2,0),(0,2,1),(0,2,2), \\
& (1,1,0),(1,1,1),(1,1,2),(1,2,0),(1,2,1),(1,2,2)\} .
\end{aligned}
$$

## Set Operations (1/7)

## كلية الحاسبات والذكاء الإصطناعي

## Union

Let $A$ and $B$ be sets. The union of the sets $A$ and $B$,
denoted by $A \cup B$, is the set that contains those
elements that are either in $A$ or in $B$, or in both.

$$
A \cup B=\{x \mid x \in A \vee x \in B\}
$$

## Set Operations (1/7)

## Union

Let $A$ and $B$ be sets. The union of the sets $A$ and $B$, denoted by $A \cup B$, is the set that contains those elements that are either in $A$ or in $B$, or in both.

$A \cup B$ is shaded.

## Set Operations (1/7)

## Union

Let $A$ and $B$ be sets. The union of the sets $A$ and $B$,
denoted by $A \cup B$, is the set that contains those
elements that are either in $A$ or in $B$, or in both.

The union of the sets $\{1,3,5\}$ and $\{1,2,3\}$ is the set $\{1,2,3,5\}$

## Set Operations (2/7)

## Intersection

Let $A$ and $B$ be sets. The intersection of the sets $A$ and
$B$, denoted by $A \cap B$, is the set that contains those elements that are in both $A$ and $B$.

$$
A \cap B=\{x \mid x \in A \wedge x \in B\}
$$

## Set Operations (2/7)

## Intersection

Let $A$ and $B$ be sets. The intersection of the sets $A$ and $B$, denoted by $A \cap B$, is the set that contains those elements that are in both $A$ and $B$.


## Set Operations (2/7)

## Intersection

Let $A$ and $B$ be sets. The intersection of the sets $A$ and
$B$, denoted by $A \cap B$, is the set that contains those elements that are in both $A$ and $B$.

The intersection of the sets $\{1,3,5\}$ and $\{1,2,3\}$ is the set $\{1,3\}$

## Set Operations (3/7)

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## Disjoint

Two sets are called disjoint if their intersection is the empty set.

$$
A \cap B=\emptyset
$$

## Set Operations (4/7)

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## Difference

Let $A$ and $B$ be sets. The difference of $A$ and $B$, denoted by $A-B$, is the set containing those elements that are in $A$ but not in $B$.

$$
A-B=\{x \mid x \in A \wedge x \notin B\}
$$

## Set Operations (4/7)

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## Difference

Let $A$ and $B$ be sets. The difference of $A$ and $B$, denoted by $A-B$, is the set containing those elements that are in $A$ but not in $B$.

$$
\begin{gathered}
A=\{1,3,5\}, \quad B=\{1,2,3\} \\
A-B=\{5\}
\end{gathered}
$$

## Set Operations (4/7)

## Difference



## Set Operations (5/7)

## Complement

Let $U$ be the universal set.
The complement of the set $A$, denoted by $\bar{A}$
An element $x$ belongs to $U$ if and only if $x \notin A$.

$$
\bar{A}=\{x \in U \mid x \notin A\}
$$

## Set Operations (5/7)

## Complement

Let $U$ be the universal set.
The complement of the set $A$, denoted by $\bar{A}$
An element $x$ belongs to $U$ if and only if $x \notin A$.

$$
\begin{gathered}
U=\{1,2,3,4,5\}, \quad A=\{1,3\} \\
\bar{A}=\{2,4,5\}
\end{gathered}
$$

## Set Operations (5/7)

## Complement


$\bar{A}$ is shaded.

## Set Operations (6/7)

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## Generalized Unions

We use the notation

$$
A_{1} \cup A_{2} \cup \cdots \cup A_{n}=\bigcup_{i=1}^{n} A_{i}
$$

to denote the union of the sets $A_{1}, A_{2}, \ldots, A_{n}$.

## Set Operations (6/7)

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## Generalized Unions


$A \cup B \cup C$ is shaded.

## Set Operations (7/7)

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## Generalized Intersections

We use the notation

$$
A_{1} \cap A_{2} \cap \cdots \cap A_{n}=\bigcap_{i=1}^{n} A_{i}
$$

to denote the intersection of the sets $A_{1}, A_{2}, \ldots, A_{n}$.

## Set Operations (7/7)

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## Generalized Intersections



## Video Lectures

All Lectures: https://www.youtube.com/playlist?list=PLxlvc-MGUsGgZIMVYOUEtUHJmfUquLiwz

https://www.youtube.com/watch?v=|FEEFRCWoBEElist=PLxlvc-

https://www.youtube.com/watch?v=RdbDHRDddn3YGlist=PLx|veMEDsEgZIMVYOEEtUHJmfUquLiwzZindex=II
https://www.youtube.com/watch?v=iSuDPBuDZzIUClist=PLxlvcMEDsEqZIMVYOEEtUHUmfUquLjwZZindex=12

Up to time 00:31:18

Up to time 00:12:46

## Thank You

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